

# **Hopf Algebras of Labeled Trees and Some Associated Differential Algebra Structures**

Robert Grossman  
University of Illinois at Chicago

Joint Work with Richard Larson  
University of Illinois at Chicago

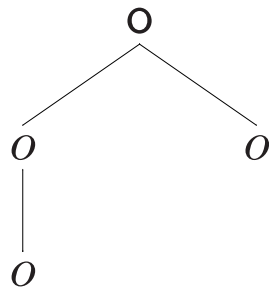
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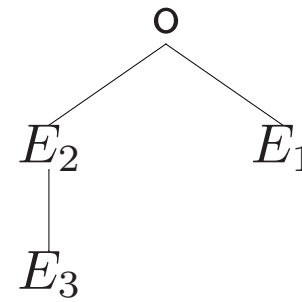
# Part 1: Introduction and Overview

# Notation

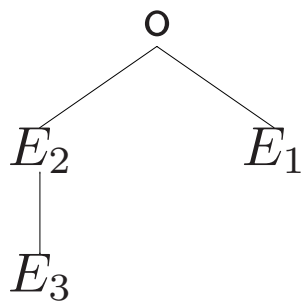
**Tree:**



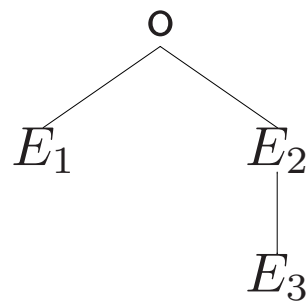
**Labeled Tree:**



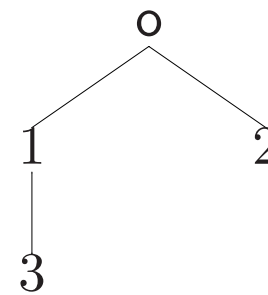
**Ordered Tree:**



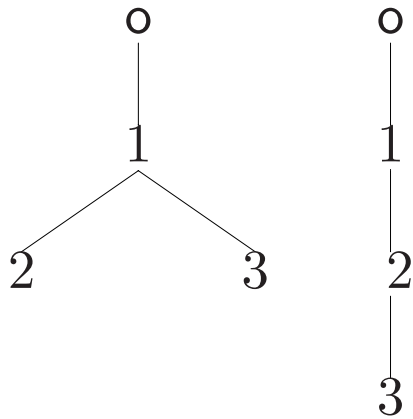
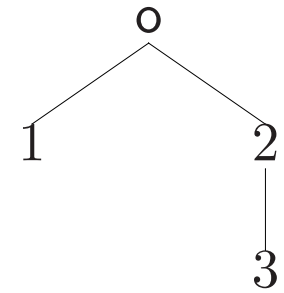
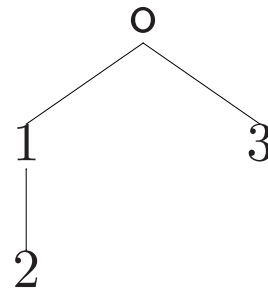
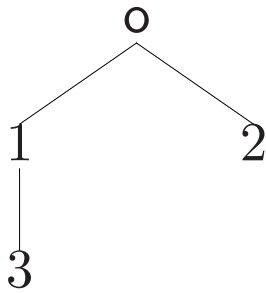
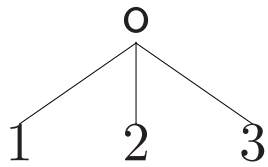
vs.



**Heap Ordered Tree:**



## Notation (con'd): Heap Ordered Trees



## Problem 1: Simplifying Higher Order Derivations

Set up: Space  $X = \mathbf{R}^n$ .  $k$ -algebra  $R$  of functions on  $X$ , say polynomials  $k[x_1, \dots, x_n]$ . Derivations  $E_1, \dots, E_M \in \text{Der}(R)$  of the form  $E_j = \sum_{\mu} a_j^{\mu} D_{\mu}$ , where  $a_j^{\mu} \in R$  and  $D_{\mu} = \partial/\partial x_{\mu}$ .

Example: Three vector fields on  $\mathbf{R}^8$ :

$$E_1 = \frac{\partial}{\partial x_1}$$

$$E_2 = \frac{\partial}{\partial x_2} - x_1 \frac{\partial}{\partial x_3} + \frac{1}{2} x_1^2 \frac{\partial}{\partial x_4} + x_1 x_2 \frac{\partial}{\partial x_5} \\ - \frac{1}{6} x_1^3 \frac{\partial}{\partial x_6} - \frac{1}{2} x_1^2 x_2 \frac{\partial}{\partial x_7} - \frac{1}{2} x_1 x_2^2 \frac{\partial}{\partial x_8}$$

$$E_3 = \frac{\partial}{\partial x_3} - x_1 \frac{\partial}{\partial x_4} - x_2 \frac{\partial}{\partial x_5} + \frac{1}{2} x_1^2 \frac{\partial}{\partial x_6} + x_1 x_2 \frac{\partial}{\partial x_7} + \frac{1}{2} x_2^2 \frac{\partial}{\partial x_8}$$

Then how do I compute efficiently  $p, q \in k\langle E_1, E_2, E_3 \rangle$ ?:

$$p = E_3E_2E_1 - E_3E_1E_2 - E_2E_1E_3 + E_1E_2E_3$$

$$q = [[E_2, E_1], E_1] = \frac{\partial}{\partial x_4} - x_1 \frac{\partial}{\partial x_6} - x_2 \frac{\partial}{\partial x_7}$$

Space $X$	Trees	Algebraic Structure	Comment
$X = \mathbf{R}^n$	Labeled Trees	bialgebra	dual is CK-Algebra
$X = G$ compact group	Labeled, Ordered Trees	bialgebra	assoc. graded dual is LR-Algebra of planar binary trees

## Problem 2: Simplifying The Action of Higher Order Derivations on Functions

Set up: Space  $X = \mathbf{R}^n$ .  $k$ -algebra  $R$  of functions on  $X$ , say polynomials  $k[x_1, \dots, x_n]$ . Derivations  $E_j \in \text{Der}(R)$ , with coefficients from  $R$ .

Example: Three vector fields on  $\mathbf{R}^8$ :

$$\begin{aligned} E_1 &= \frac{\partial}{\partial x_1} \\ E_2 &= \frac{\partial}{\partial x_2} - x_1 \frac{\partial}{\partial x_3} + \frac{1}{2} x_1^2 \frac{\partial}{\partial x_4} + x_1 x_2 \frac{\partial}{\partial x_5} \\ &\quad - \frac{1}{6} x_1^3 \frac{\partial}{\partial x_6} - \frac{1}{2} x_1^2 x_2 \frac{\partial}{\partial x_7} - \frac{1}{2} x_1 x_2^2 \frac{\partial}{\partial x_8} \end{aligned}$$

Then how do I compute efficiently  $p(r)$ , where  $p \in k\langle E_1, E_2, E_3 \rangle$  and  $r \in R$ ?

$$(E_2 E_2 E_1 + E_1)(3x_1^3 x_2 - 2x_1 x_7) \quad [E_2, E_1](x_3) = 1, \quad [[E_2, E_1], E_1](x_4) = 1,$$

<b>Space <math>X</math></b>	<b>Trees</b>	<b>Algebraic Structure</b>	<b>Comment</b>
$X = \mathbf{R}^n$	Labeled Trees	H-Module Algebra	H = bialgebra of trees
$X = G$ compact group	Labeled, Ordered Trees	$R/k$ -bialgebra	requires connection on G

## $R/k$ -bialgebras

$R/k$ -bialgebras encode differential algebra structure and were introduced by Nichols and Weisfeiler in 1982:

W. Nichols and B. Weisfeiler, Differential formal groups of J. F. Ritt, Amer. J. Math. 104 (1982) 943-1003.

W. Nichols, The Kostant structure theorems for  $K/k$ -Hopf algebras, Journal Algebra 97 (1985) 313-328.

### Problem 3: Computing Flows and their Symmetries

Set up: Space  $X = G$ , a group. Fix a basis  $\{Y_1, \dots, Y_M\}$  of the Lie algebra. Fix a derivation  $E = \sum_{\mu} a^{\mu} Y_{\mu}$  with coefficients from the coordinate ring  $R$ .

Example: Compute symbolic approximations to the flow  $\exp tE$  or

$$\exp\left(t \begin{array}{c} \circ \\ | \\ E \end{array} \right) = 1 + t \begin{array}{c} \circ \\ | \\ E \end{array} + \frac{t^2}{2!} \begin{array}{c} \circ \\ | \\ E \end{array} \cdot \begin{array}{c} \circ \\ | \\ E \end{array} + \dots$$

Space $X$	Trees	Algebraic Structure	Comment
$X = G$	Heap Ordered Trees	bialgebra	assoc. graded dual is MR-Hopf algebra of permutations

## Cayley 1857

Cayley noted that repeated differentiation led to trees and showed the following:

Let  $t_n$  be the number of trees with  $n$  nodes. Cayley showed that if  $T(z) = \sum_{n=1}^{\infty} t_n z^n$ , then

$$T(z) = z \prod_{n=1}^{\infty} (1 - z^n)^{-t_n}.$$

A. Cayley, On the theory of the analytical forms called trees, in “Collected Mathematical Papers of Arthur Cayley,” Cambridge University Press, Cambridge, 1890, 3, 242-246.

## Part 2: Simplifying Higher Order Derivations in $\mathbb{R}^n$

This is Problem 1 in the commutative case.

## Example: Simplifying Higher Order Derivations

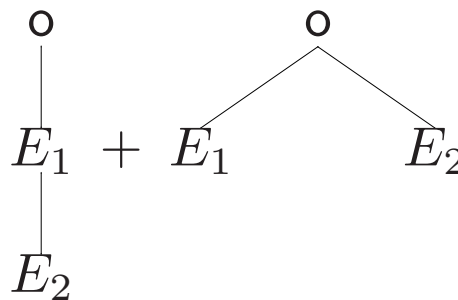
Consider three vector fields

$$E_1 = a_1(x)D_1 + \cdots + a_N(x)D_N, \quad E_2 = b_1(x)D_1 + \cdots + b_N(x)D_N,$$

$$E_3 = c_1(x)D_1 + \cdots + c_N(x)D_N.$$

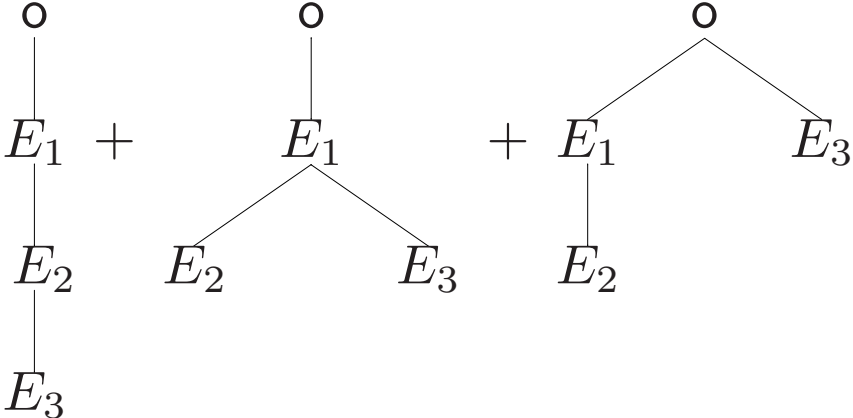
where  $D_i = \partial/\partial x_i$ , and  $a_i(x)$ ,  $b_i(x)$  and  $c_i(x)$  are smooth functions on  $\mathbf{R}^N$ . Now

$$E_2 \cdot E_1 = \sum b_j(D_j a_i)D_i + \sum b_j a_i D_j D_i$$

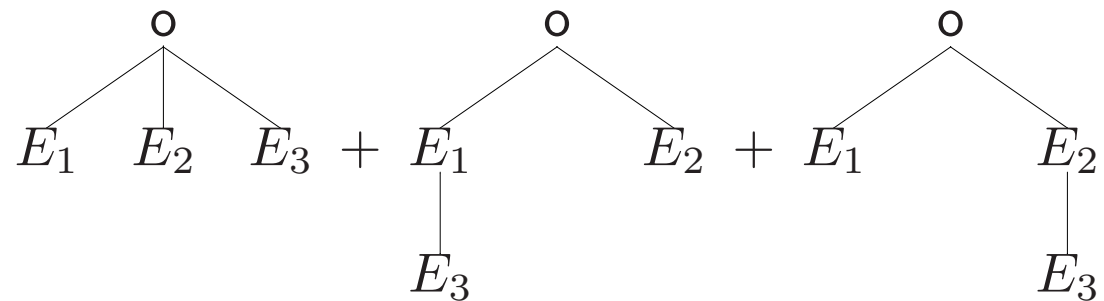


Now  $E_3 \cdot E_2 \cdot E_1$  is equal to

$$\sum c_k (D_k b_j) (D_j a_i) D_i + \sum c_k b_j (D_k D_j a_i) D_i + \sum c_k b_j (D_j a_i) D_k D_i$$



$$+ \sum c_k b_j a_i D_k D_j D_i + \sum c_k b_j (D_k a_i) D_j D_i + \sum c_k (D_k b_j) a_i D_j D_i.$$



## Theorem (1989).

If a family of rooted trees (possibly with additional structure) and four operations (DeleteRoot, Nodes, Restriction, AttachRoot) satisfies 9 axioms, then the family forms a Hopf algebra.

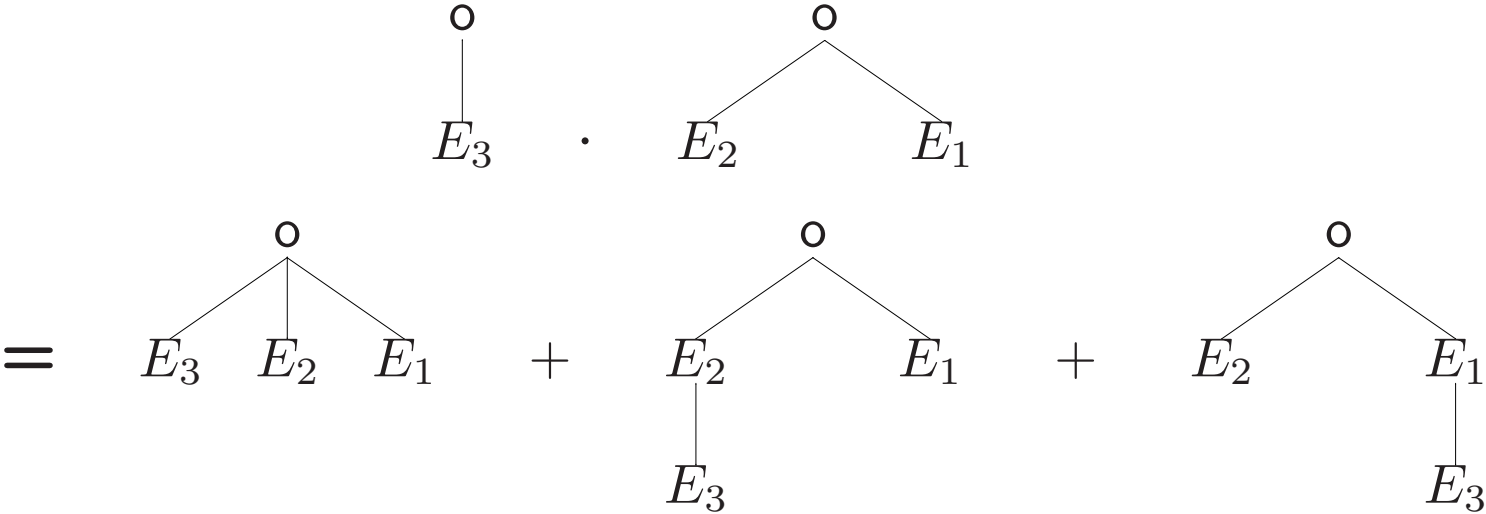
### Examples:

- trees
- labeled trees
- ordered trees
- heap ordered trees

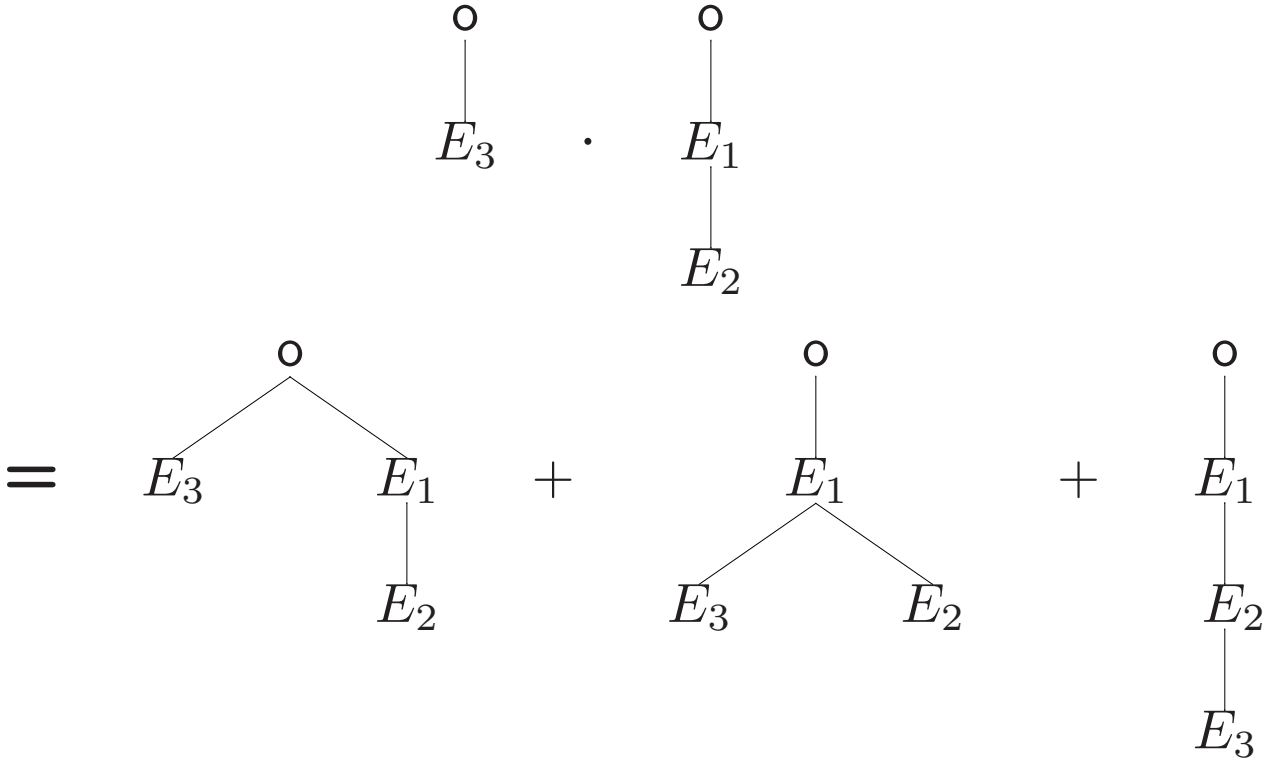
## Product - Example

$$\begin{array}{c} \circ \\ | \\ E_2 \end{array} \cdot \begin{array}{c} \circ \\ | \\ E_1 \end{array} = \begin{array}{c} \circ \\ | \\ E_1 \\ | \\ E_2 \end{array} + \begin{array}{c} \circ \\ / \quad \backslash \\ E_1 \quad E_2 \end{array}$$

# Product - Example



# Product - Example



# Coproduct

$$\Delta\left(\begin{array}{c} \circ \\ | \\ E_1 \end{array}\right) = \mathbf{1} \otimes \begin{array}{c} \circ \\ | \\ E_1 \end{array} + \begin{array}{c} \circ \\ | \\ E_1 \end{array} \otimes \mathbf{1}$$

## Coproduct (cont'd)

$$\begin{aligned}
 \Delta( \begin{array}{c} \circ \\ \diagdown \quad \diagup \\ E_1 \quad E_2 \end{array} ) &= \mathbf{1} \otimes \begin{array}{c} \circ \\ \diagdown \quad \diagup \\ E_1 \quad E_2 \end{array} \\
 &+ \begin{array}{c} \circ \\ | \\ E_1 \end{array} \otimes \begin{array}{c} \circ \\ | \\ E_2 \end{array} + \begin{array}{c} \circ \\ | \\ E_2 \end{array} \otimes \begin{array}{c} \circ \\ | \\ E_1 \end{array} \\
 &+ \begin{array}{c} \circ \\ \diagdown \quad \diagup \\ E_1 \quad E_2 \end{array} \otimes \mathbf{1}
 \end{aligned}$$

## Example:

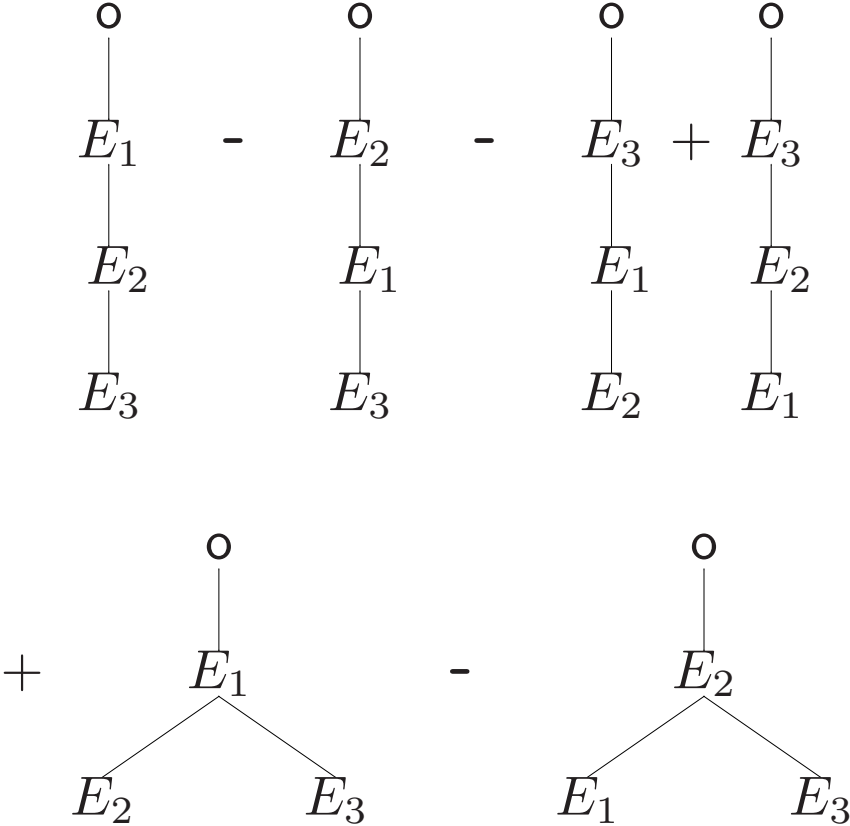
An expression such as

$$p = E_3 E_2 E_1 - E_3 E_1 E_2 - E_2 E_1 E_3 + E_1 E_2 E_3$$

corresponds to 24 trees, 18 of which cancel (each cancellation saves  $O(N^3)$  differentiations).

$$\begin{aligned} & \sum a_3^{\mu_3} (D_{\mu_3} a_2^{\mu_2}) (D_{\mu_2} a_1^{\mu_1}) D_{\mu_1} - \sum a_3^{\mu_3} (D_{\mu_3} a_1^{\mu_2}) (D_{\mu_2} a_2^{\mu_1}) D_{\mu_1} \\ & - \sum a_2^{\mu_3} (D_{\mu_3} a_1^{\mu_2}) (D_{\mu_2} a_3^{\mu_1}) D_{\mu_1} + \sum a_1^{\mu_3} (D_{\mu_3} a_2^{\mu_2}) (D_{\mu_2} a_3^{\mu_1}) D_{\mu_1} \\ & + \sum a_3^{\mu_3} a_2^{\mu_2} (D_{\mu_3} D_{\mu_2} a_1^{\mu_1}) D_{\mu_1} - \sum a_3^{\mu_3} a_1^{\mu_2} (D_{\mu_3} D_{\mu_2} a_2^{\mu_1}) D_{\mu_1}, \end{aligned}$$

The surviving six trees are:



# Summary

Space $X$	Trees	Algebraic Structure	Comment
$X = \mathbf{R}^n$	Labeled Trees	bialgebra	helpful with simplification of expressions involving higher order derivations (each tree that cancels results in $O(N^r)$ fewer computations)

## **Part 3: Simplifying The Action of Higher Order Derivations on Functions**

This is Problem 2 in the commutative case.

## Recall the Co-Product

$$\Delta\left(\begin{array}{c} \circ \\ | \\ E \end{array}\right) = \mathbf{1} \otimes \begin{array}{c} \circ \\ | \\ E \end{array} + \begin{array}{c} \circ \\ | \\ E \end{array} \otimes \mathbf{1}$$

We want:

$$\begin{array}{c} \circ \\ | \\ E \end{array} (f) \leftrightarrow E(f),$$

where  $E \in \text{Der}(R)$  and  $f \in R$ .

## Example

Analysis of actions of higher order derivations on algebras of functions using  $H$ -module algebras.

$$\begin{aligned}
 \underset{E}{\overset{\circ}{\mid}} (\mathbf{fg}) &= \mathbf{f} \underset{E}{\overset{\circ}{\mid}} \mathbf{g} + \mathbf{g} \underset{E}{\overset{\circ}{\mid}} \mathbf{f} \\
 &= (\mathbf{1} \otimes \underset{E}{\overset{\circ}{\mid}} + \underset{E}{\overset{\circ}{\mid}} \otimes \mathbf{1})(\mathbf{f} \otimes \mathbf{g}) \\
 &= \Delta(\underset{E}{\overset{\circ}{\mid}})(\mathbf{f} \otimes \mathbf{g})
 \end{aligned}$$

## H-Module Algebras

Let  $R$  be a commutative  $k$ -algebra, and let  $H$  be a  $k$ -bialgebra. The algebra  $R$  is a *left  $H$ -module algebra* if  $R$  is a left  $H$ -module for which

$$h \cdot (ab) = \sum_{(h)} (h_{(1)} \cdot a)(h_{(2)} \cdot b),$$

where  $h \in H$ ,  $\Delta(h) = \sum_{(h)} h_{(1)} \otimes h_{(2)}$ , and  $a, b \in R$ .

**Why is this important?** Many practical computations involve analyzing the actions of differential operators on specific functions. This translates into understanding the actions of specific bialgebras algebras of trees on specific algebras of functions or  $H$ -module algebras.

## Theorem (1990).

- Fix derivations  $E_j = \sum_{\mu} a_j^{\mu} D_{\mu}$ , with  $a_j^{\mu} \in R$ ,  $D_{\mu} = \frac{\partial}{\partial x_{\mu}}$ .
- Let  $k\langle E_1, \dots, E_M \rangle$  be the free associative algebra generated by the formal symbols  $E_1, \dots, E_M$ . Let  $H = k\{\mathcal{LT}(E_1, \dots, E_M)\}$  be the Hopf algebra consisting of rooted trees labeled with derivations  $E_j$ . Let  $\text{Diff}(E_1, \dots, E_M)$  be the higher order derivations generated by the derivations  $E_1, \dots, E_M$ .

Then one can define a morphism  $\psi$  and an action so that  $R$  is an  $H$ -module and the following diagram commutes:

$$\begin{array}{ccc}
 k\langle E_1, \dots, E_M \rangle & \rightarrow & k\{\mathcal{LT}(E_1, \dots, E_M)\} \\
 \searrow & & \downarrow \psi \\
 & & \text{Diff}(E_1, \dots, E_M)
 \end{array} \tag{1}$$



## $\psi$ map in general:

Let  $h \in k\{\mathcal{LT}(E_1, \dots, E_M)\}$  be a labeled tree. Number the root of  $h$  with 0, and number the other nodes  $1, \dots, m$ . Suppose that node  $i$ ,  $i \geq 0$ , has children  $j_1, \dots, j_k$ . Define

$$R(i) = \begin{cases} D_{\mu_{j_k}} \cdots D_{\mu_{j_1}} f & \text{if } i = 0, \\ D_{\mu_{j_k}} \cdots D_{\mu_{j_1}} a_i^{\mu_i} & \text{otherwise.} \end{cases}$$

Define

$$h \cdot f = \sum_{\mu_1, \dots, \mu_m=1}^N R(m) \cdots R(1)R(0),$$

for  $f \in R$ . This makes  $R$  into a left  $H$ -module algebra.

## Theorem (1990)

For all functions  $a, b \in R$  and for all labeled trees  $h \in k\{\mathcal{LT}(E_1, \dots, E_M)\}$ ,

$$((\psi \otimes \psi) \cdot \Delta(h))(a \otimes b) = \psi(h)(ab)$$

**We conclude:** The  $\psi$  map defines an  $H$ -module algebra structure on  $R$  which makes the diagram commute.

# Summary

<b>Space <math>X</math></b>	<b>Trees</b>	<b>Algebraic Structure</b>	<b>Comment</b>
$X = \mathbf{R}^n$	Labeled Trees	H-Module Algebra	helpful when simplifying action of higher order derivations on functions

## Question

- What happens when  $X$  is not  $\mathbf{R}^N$  and the  $D_\mu$  defining the derivations

$$E_j = \sum_{\mu} a_j^{\mu}(x) D_{\mu}$$

do not commute?

## Part 4: Simplifying Higher Order Derivations for $\text{Diff}(G)$

This is Problem 1 in the non-commutative case.

## Key Observation - Defining natural $H$ -modules in the non-Abelian case requires using a connection.

Let  $R$  be a commutative  $k$ -algebra, and let  $\mathbf{D}$  be a Lie algebra of derivations of  $R$ . A *connection* is a map  $\mathbf{D} \times \mathbf{D} \rightarrow \mathbf{D}$  sending  $(E, F) \in \mathbf{D} \times \mathbf{D}$  to  $\nabla_E F \in \mathbf{D}$  satisfying

- $\nabla_{E_1+E_2} F = \nabla_{E_1} F + \nabla_{E_2} F$
- $\nabla_E (F_1 + F_2) = \nabla_E F_1 + \nabla_E F_2$
- $\nabla_{f \cdot E} F = f \cdot \nabla_E F$
- $\nabla_E (f \cdot F) = f \cdot \nabla_E F + E(f)F$

where  $E, F \in \mathbf{D}$ ,  $f \in R$ .

**We use this connection as follows:**

$$\begin{array}{c} \circ \\ | \\ E \end{array} (r) = E(r)$$

$$\begin{array}{c} \circ \\ | \\ E \\ | \\ F \end{array} (r) = \nabla_F E(r)$$

and extend using induction and a consistency requirement to larger trees.

## Theorem (2005)

Let  $R$  be a commutative  $k$ -algebra, and let  $\nabla_E F$  be a connection on the Lie algebra  $\mathbf{D}$  of derivations of  $R$ . Then the construction above gives a  $k\{\mathcal{L}\mathcal{T}(\mathbf{D})\}$ -module structure on  $R$ . This module structure induces a map  $\psi : k\{\mathcal{L}\mathcal{T}(\mathbf{D})\} \longrightarrow \text{End}(R)$ . so that the following diagram commutes:

$$\begin{array}{ccc} k\langle E_1, \dots, E_M \rangle & \rightarrow & k\{\mathcal{L}\mathcal{T}(\mathbf{D})\} \\ & \searrow & \downarrow \psi \\ & & \text{Diff}(R) \subset \text{End}(R) \end{array} \quad (2)$$

## $R/k$ -Algebras

Let  $R$  be a  $k$ -algebra. A  $R/k$ -algebra is a  $k$ -algebra  $B$  into which  $R$  is embedded. Note that this makes  $B$  into a left and right  $R$ -module, and that  $(rb)s = r(bs)$  for all  $r, s \in R, b \in B$ , and that also  $(rb)c = r(bc)$ ,  $(br)c = b(rc)$ , and  $(bc)r = b(cr)$  for all  $b, c \in B, r \in R$ .

Example:  $\text{Diff}(\text{Der}(R))$  is a  $R/k$ -algebra.

**It turns out that there is a version of the construction above using  $R/k$ -algebras and  $R/k$ -bialgebras.**

## $R/k$ -bialgebra

A  $R/k$ -bialgebra is a  $R/k$ -algebra  $B$  together with  $R$ -module maps  $\Delta : B \rightarrow B \otimes_R B$  and  $\epsilon : B \rightarrow R$  satisfying

- $B$  together with the maps  $\Delta$  and  $\epsilon$  is a coalgebra over  $R$ .
- $\Delta(1) = 1 \otimes 1$ .
- For all  $b, c \in B$ , if  $\Delta(b) = \sum_i b_i \otimes b'_i$  and  $\Delta(c) = \sum_j c_j \otimes c'_j$  are any representations of  $\Delta(b), \Delta(c) \in B \otimes_R B$ , then  $\Delta(bc) = \sum_{i,j} b_i c_j \otimes b'_i c'_j$ .
- $\epsilon(1) = 1$ .
- $\epsilon(bc) = \epsilon(b\epsilon(c))$ .

**Theorem.**  $R \otimes_k k\{\mathcal{LT}(\mathbf{Der}(R))\}$  is a  $R/k$ -bialgebra.

Let  $R$  be a commutative  $k$ -algebra, let  $\mathbf{Der}(R)$  denote derivations of  $R$ , and let  $\mathbf{Diff}(\mathbf{Der}(R))$  denote differential operators generated by  $\mathbf{Der}(R)$ .

- The tensor product  $R \otimes_k \mathbf{Diff}(\mathbf{Der}(R))$  is  $R/k$  – bialgebra.
- Let  $k\{\mathcal{LT}(\mathbf{Der}(R))\}$  denote the bialgebra of trees labeled with elements of  $\mathbf{Der}(R)$ . Then the tensor product  $R \otimes_k k\{\mathcal{LT}(\mathbf{Der}(R))\}$  is a  $R/k$ -bialgebra.
- The map

$$R \otimes_k k\{\mathcal{LT}(\mathbf{Der}(R))\} \rightarrow R \otimes_k \mathbf{Diff}(\mathbf{Der}(R))$$

is a  $R/k$ -bialgebra homomorphism.

# Summary

<b>Space <math>X</math></b>	<b>Trees</b>	<b>Algebraic Structure</b>	<b>Comment</b>
$X = G$ compact group	Labeled, Ordered Trees	$R/k$ - bialgebra	This structures helps with simplifying higher order derivations on groups

## Part 5: Computing Flows on Groups $G$

This is Problem 3 in the non-commutative case.

# Computing geometrically stable numerical algorithms.

Goal: Develop a numerical algorithm to approximate flow on a Lie Group  $G$  of

$$\dot{x}(t) = E(x(t)), \quad x(0) = x^0 \in G$$

- $E = \sum_{\mu} a^{\mu}(x) Y_{\mu}$
- $a^{\mu}(x)$  in coordinate ring  $R$
- $Y_{\mu}$  basis Lie algebra  $\mathfrak{g}$

## Classical Runge Kutta Algorithms:

Consider the differential equation:  $x'(t) = E(x(t))$ , subject to the initial condition  $x(0) = x^0$ . Find coefficients  $c_{i,j}$  and  $b_i$

$$X_1 = x_0$$

$$X_2 = x_0 + hc_{2,1}E(X_1)$$

$$X_3 = x_0 + hc_{3,1}E(X_1) + hc_{3,2}E(X_2)$$

⋮

$$X_r = x_0 + hc_{r,1}E(X_1) + hc_{r,2}E(X_2) + \cdots + hc_{r,r-1}E(X_{r-1})$$

and let

$$x_1 = x_0 + \sum_{i=1}^r hb_i E(X_i).$$

so that

$$x_1 - x(h) = O(h^k)$$

## Runge Kutta Type Algorithms on Lie Groups:

Fix a point  $p \in G$  and consider the ansatz:

$$\bar{E}_1 = \sum a^\mu(p) Y_\mu \in \mathfrak{g}$$

$$\bar{E}_2 = \sum a^\mu(\exp(hc_{21}\bar{E}_1 \cdot p)) Y_\mu \in \mathfrak{g}$$

$$\bar{E}_3 = \sum a^\mu(\exp(hc_{32}\bar{E}_2 \exp(hc_{31}\bar{E}_1 \cdot p))) Y_\mu \in \mathfrak{g}$$

**Compute:**

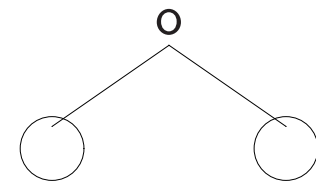
$$LHS = \exp hc_3 \bar{E}_3 \exp hc_2 \bar{E}_2 \exp hc_1 \bar{E}_1 - \exp hE = 0 + O(h^4)$$

$$\begin{aligned}
LHS &= hc_1E + hc_2E + hc_3E \\
&+ h^2c_2c_{21}DE \cdot E + h^2c_3c_{31}DE \cdot E + h^2c_3c_{32}DE \cdot E \\
&\quad + h^3c_3c_{32}c_{21}DE \cdot DE \cdot E \\
&\quad + h^3c_2\frac{c_{21}^2}{2!}D^2E(E, E) + h^3c_3\frac{c_{31}^2}{2!}D^2E(E, E) \\
&\quad + h^3c_3\frac{c_{32}^2}{2!}D^2E(E, E) + h^3c_3c_{31}c_{32}D^2E(E, E) \\
&- hE - \frac{h^2}{2!}DE \cdot E - \frac{h^3}{3!}D^2E(E, E) - \frac{h^3}{3!}DE \cdot DE \cdot E
\end{aligned}$$

$$LHS = hc_1E + hc_2E + hc_3E$$



$$+h^2c_2c_{21}DE \cdot E + h^2c_3c_{31}DE \cdot E + h^2c_3c_{32}DE \cdot E$$

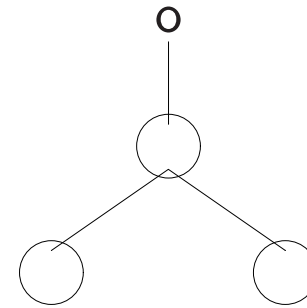


$$+h^3 c_3 c_{32} c_{21} DE \cdot DE \cdot E$$



$$+h^3 c_2 \frac{c_{21}^2}{2!} D^2 E(E, E) + h^3 c_3 \frac{c_{31}^2}{2!} D^2 E(E, E)$$

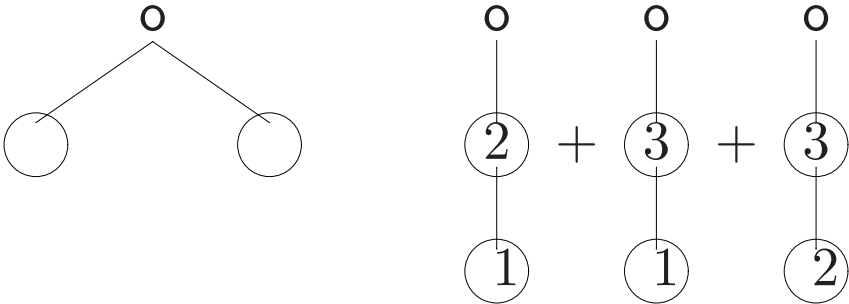
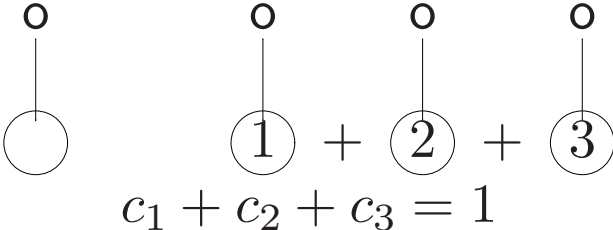
$$h^3 c_3 \frac{c_{32}^2}{2!} D^2 E(E, E) + h^3 c_3 c_{31} c_{32} D^2 E(E, E)$$



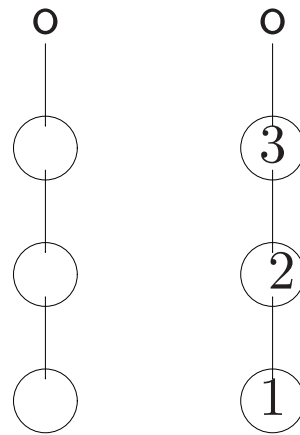
$$-hE - \frac{h^2}{2!} DE \cdot E - \frac{h^3}{3!} D^2 E(E, E) - \frac{h^3}{3!} DE \cdot DE \cdot E$$

$$+O(h^4)$$

We have:



$$c_2c_{21} + c_3c_{31} + c_3c_{32} = \frac{1}{2!}$$



$$c_3 + c_{32} + c_{21} = \frac{1}{3!}$$

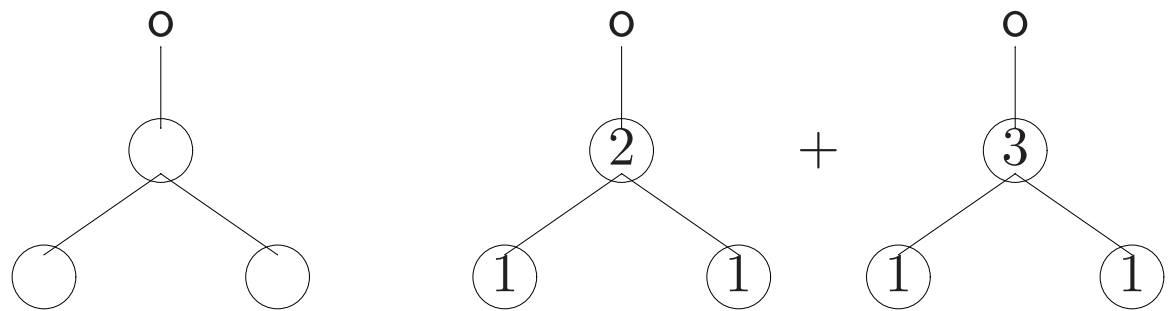


Diagram illustrating two tree structures:

- Tree 4: Root node 0, child node 3, two children 2.
- Tree 5: Root node 0, child node 3, children 1 and 2.

A plus sign (+) is placed between the two trees.

$$\frac{c_2 c_{21}^2}{2!} + \frac{c_3 c_{31}^2}{2!} + \frac{c_3 c_{32}^2}{2!} + c_3 c_{31} c_{32} = \frac{1}{3!}$$

## Theorem (1993)

Third order RK algorithms on Lie Groups arise by solving the constraint equations:

$$1. c_1 + c_2 + c_3 = 1$$

$$2. c_2c_{21} + c_3c_{31} + c_3c_{32} = 1$$

$$3. \frac{c_2c_{21}^2}{2!} + \frac{c_3}{2!}(c_{31} + c_{32})^2 = 1/3!$$

$$4. c_3c_{32}c_{21} = 1/3!$$

$$5. 3(c_2^2c_{21} + c_3^2(c_{32} + c_{31})) + 6c_2c_3c_{21} = 1$$

# Summary

<b>Space <math>X</math></b>	<b>Trees</b>	<b>Algebraic Structure</b>	<b>Comment</b>
$X = G$	RK- Ordered Trees	bialgebra	This codes RK- algorithms on groups.

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**Thank you.**